

Befehlssatz des D-CORE Prozessors

Mnemonic	Codierung	Hex	Bedeutung
ALU-Operationen			
mov	0010 0000 <yyyy> <xxxx>	20yx	$R[x] = R[y]$
addu	0010 0001 <yyyy> <xxxx>	21yx	$R[x] = R[x] + R[y]$
addc	0010 0010 <yyyy> <xxxx>	22yx	$R[x] = R[x] + R[y] + C$ (updates C)
subu	0010 0011 <yyyy> <xxxx>	23yx	$R[x] = R[x] - R[y]$
and	0010 0100 <yyyy> <xxxx>	24yx	$R[x] = R[x] \text{ AND } R[y]$ (bitwise)
or	0010 0101 <yyyy> <xxxx>	25yx	$R[x] = R[x] \text{ OR } R[y]$ (bitwise)
xor	0010 0110 <yyyy> <xxxx>	26yx	$R[x] = R[x] \text{ XOR } R[y]$ (bitwise)
not	0010 0111 <****> <xxxx>	27*x	$R[x] = \text{NOT } R[x]$ (bitwise)
Shift-Operationen			
lsl	0010 1000 <yyyy> <xxxx>	28yx	$R[x] = R[x] \ll R[y].<3:0>$
lsr	0010 1001 <yyyy> <xxxx>	29yx	$R[x] = R[x] \gg R[y].<3:0>$
asr	0010 1010 <yyyy> <xxxx>	2Ayx	$R[x] = R[x] \gg R[y].<3:0>$
lslc	0010 1100 <****> <xxxx>	2C*x	$R[x] = R[x] \ll 1, C=R[X].15$
lsrc	0010 1101 <****> <xxxx>	2D*x	$R[x] = R[x] \gg 1, C=R[X].0$
asrc	0010 1110 <****> <xxxx>	2E*x	$R[x] = R[x] \gg 1, C=R[X].0$
Vergleichs-Operationen			
cmpe	0011 0000 <yyyy> <xxxx>	30yx	$C = (R[x] == R[y])$
cmpne	0011 0001 <yyyy> <xxxx>	31yx	$C = (R[x] != R[y])$
cmpgt	0011 0010 <yyyy> <xxxx>	32yx	$C = (R[x] > R[y])$ (signed)
cmplt	0011 0011 <yyyy> <xxxx>	33yx	$C = (R[x] < R[y])$ (signed)
Immediate-Operationen			
movi	0011 0100 <cccc> <xxxx>	34cx	$R[x] = 0x000c$
addi	0011 0101 <cccc> <xxxx>	35cx	$R[x] = R[x] + 0x000 \langle c \rangle$
subi	0011 0110 <cccc> <xxxx>	36cx	$R[x] = R[x] - 0x000 \langle c \rangle$
andi	0011 0111 <cccc> <xxxx>	37cx	$R[x] = R[x] \text{ AND } 0x000 \langle c \rangle$
lsli	0011 1000 <cccc> <xxxx>	38cx	$R[x] = R[x] \ll \langle c \rangle$
lsri	0011 1001 <cccc> <xxxx>	39cx	$R[x] = R[x] \gg \langle c \rangle$
bseti	0011 1010 <cccc> <xxxx>	3Acx	$R[x] = R[x] (1 \ll \langle c \rangle)$ (set bit)
bclri	0011 1011 <cccc> <xxxx>	3Bcx	$R[x] = R[x] \& !(1 \ll \langle c \rangle)$ (clear bit)
Speicher-Operationen			
ldw	0100 <cccc> <yyyy> <xxxx>	4cyx	$R[x] = \text{MEM}(R[y] + (0x000 \langle c \rangle \ll 1))$
stw	0101 <cccc> <yyyy> <xxxx>	5cyx	$\text{MEM}(R[y] + (0x000 \langle c \rangle \ll 1)) = R[x]$
Kontrollfluss			
br	1000 <iiii> <iiii> <iiii>	8iii	$\text{PC} = \text{PC} + 2 + \langle imm12 \rangle$
jsr	1001 <iiii> <iiii> <iiii>	9iii	$R[15] = \text{PC} + 2; \text{PC} = \text{PC} + 2 + \langle imm12 \rangle$ (call)
bt	1010 <iiii> <iiii> <iiii>	Aiii	$(C=1) ? \text{PC} = \text{PC} + 2 + \langle imm12 \rangle : \text{PC} = \text{PC} + 2$
bf	1011 <iiii> <iiii> <iiii>	Biii	$(C=0) ? \text{PC} = \text{PC} + 2 + \langle imm12 \rangle : \text{PC} = \text{PC} + 2$
jmp	1100 <****> <****> <xxxx>	C*x	$\text{PC} = R[x]$
halt	1111 <****> <****> <****>	F***	halt andere Opcodes illegal

<xxxx>, x: 4-bit Index des Quell- und Zielregisters RX <xxxx> – binär, x – hex
 <yyyy>, y: 4-bit Index des Quellregisters RY <yyyy> – binär, y – hex
 <cccc>, c: 4-bit Konstante IMM4 <cccc> – binär, c – hex
 iii: 12-bit sign-extended Konstante IMM12 <iiii> – binär, i – hex
 <****>, *: don't care